

Quantum Computation and Information

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Source: Nielsen and Chuang: Quantum Computation and Quantum Information (2000)

Definition of quantum computing: study of information processing tasks which use quantum mechanical systems

Combination of quantum mechanics, computer science, information theory

Quantum mechanics: most successful theory to describe microscopic world, but is often counterintuitive (developed in 1920s)
certain [↓] operations which are trivial in classical information are not even possible in quantum case

example: cloning of data

only imperfect cloning is possible

- complete control over quantum systems (single)
- recent developments: single atom traps
 - can observe properties of single atom isolated from surroundings
- can we do quantum information processing?
 - state-of-the-art: manipulation of a few qubits
 - some success in quantum cryptography

Computer Science

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Turing - developed abstract notion of what we would call a computer (1936)

Universal Turing machine: completely captures the performance of algorithms in principle

- first computers:

- Z3 (Konrad Zuse, 1941)

- 1947 Bardeen, Brattain, Shockley \Rightarrow transistor

- Moore's law: computer power for a given cost doubles once every two years

(approximately true since 1960's

up to now)

- we expect that Moore's law will not hold in the near

future \Rightarrow size of electronic devices are smaller and smaller \Rightarrow are approaching the molecular/atomic length scale where quantum mechanics is important

to maintain Moore's law \Rightarrow need quantum computing

(or some other paradigm change)

- quantum computer - can be simulated

using a classical computer \Rightarrow but not efficiently

efficient: means that effort needed

scales as polynomial of the size

of the problem

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Church - Turing hypothesis: Any algorithmic process can be simulated efficiently using a Turing machine

- hypothesis challenged by randomized algorithms

1970s Solovay - Strassen: can determine whether a number is prime or not within some probability

↓ ↓
uses random number generator

→ extend Church - Turing hypothesis to probabilistic case

Deutsch: can one derive a Church - Turing hypothesis based on the laws of physics?

- can one construct a device to simulate an arbitrary physical system?

↓
considered quantum analogs of devices in Turing machines

↓
~~but~~ hypothesis is still an open issue
but Church - Turing hypothesis has been challenged → for some cases quantum version of algorithm is more efficient

- other examples: Shor's - discrete logarithm algorithm
- finding prime factors of an integer

quantum computers seem to be more efficient than Turing machines

Grover: quantum search algorithms

Feynmann: simulation of a quantum system
upto now an impossible task for systems
of meaningful size

- issue with quantum computing: our intuition is rooted in the classical world
→ hard to construct quantum algorithms

information theory

- aim: quantity resources for sending information from sender to receiver

- issues: noisy channel, error-correction

- Shannon: theorems which accomplish this

- for quantum information theory
no analogous theorem

- nevertheless: in principle possible to send information even in the presence of noise, error-correcting also possible

- cryptography: transmit message which cannot be read by third party

→

- quantum cryptography:

- measurement in quantum mechanics changes state
→ third party interference detectable

Quantum Bits

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bit - fundamental concept of classical computation \Rightarrow 0, 1

qubit - same for quantum computing

- for now introduce conceptually \Rightarrow they can also be realized

- classical bit: state $|0\rangle$ or $|1\rangle$

quantum bit: in superposition state

$$|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$$

α, β - complex numbers

$|0\rangle, |1\rangle$ - known as basis states

orthonormal: $\langle 0|0\rangle = \langle 1|1\rangle = 1$

$$\langle 0|1\rangle = \langle 1|0\rangle = 0$$

- classical bit \Rightarrow can be determined
state is either 1 or 0

- quantum bit \Rightarrow state can not be determined

measurement gives 0 with probability $|\alpha|^2$

or

1 with probability $|\beta|^2$

$$|\alpha|^2 + |\beta|^2 = 1 \Rightarrow \text{state of qubit is}$$

a unit vector in complex space

- classical state observable / quantum state unobservable

\Downarrow

heart of quantum computation!!!

- even though α, β cannot be measured directly, (6) manipulations on qubits (gates) are possible



consequences experimentally verifiable

- classical bit: exists in state 0 or 1
- qubit: exists in superposition state until observed!

example: $\frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle$

before measurement: 50% 0 50% 1 after measurement: 0 or 1 (100%)

- qubits have a number of experimental realizations
 - Stern-Gerlach → spin of individual atoms
 - ground/excited state of atom/molecule

- geometric representation of qubit

$$|\psi\rangle = e^{i\phi} \left(\cos\frac{\theta}{2} |0\rangle + e^{i\varphi} \sin\frac{\theta}{2} |1\rangle \right)$$

α, β - 4 unknowns → 3 after normalization

θ, ϕ, φ - 3 unknowns

- turns out ϕ is (usually) inconsequential

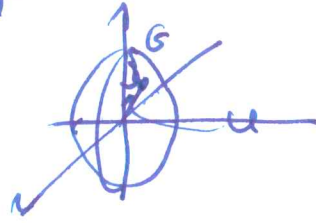
θ, φ - point on unit 3D sphere



$$|\psi\rangle = \cos\frac{\theta}{2} |0\rangle + e^{i\varphi} \sin\frac{\theta}{2} |1\rangle$$

Bloch sphere: visualization of state of single qubit

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gates (operations on qubits) \Rightarrow operations on unit vector \Rightarrow rotations

information stored in qubit?

- in principle: surface of sphere $\Rightarrow \infty$ - points
- problem: can not easily be recovered
- measurement/observation: qubit falls into specific state $|0\rangle$ or $|1\rangle$ and qubit stays there ~~ever after~~ (until next nontrivial operation)
- "collapse" \rightarrow one of fundamental postulates of quantum mechanics
- to reconstruct state $\alpha/\beta \Rightarrow$ need an ensemble of identically prepared states with an infinite number of members

- if we do not measure qubit:

remains in state α, β until further operations:

$$\text{then } \alpha\beta \rightarrow \alpha'\beta' \rightarrow \alpha''\beta''$$

according to propagation

$\Rightarrow \alpha, \beta$ - coefficients of hidden information

Multiple qubits

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- two qubits: four states: $|00\rangle, |01\rangle, |10\rangle, |11\rangle$

$$|\varphi\rangle = \alpha_{00}|00\rangle + \alpha_{01}|01\rangle + \alpha_{10}|10\rangle + \alpha_{11}|11\rangle$$

- One can measure 1st qubit, 2nd qubit or both simultaneously

- $|\alpha\rangle$ $\begin{matrix} \swarrow & \text{second qubit} \\ \uparrow & \text{1st qubit} \end{matrix}$

$$|\alpha_{00}|^2 + |\alpha_{01}|^2 + |\alpha_{10}|^2 + |\alpha_{11}|^2 = 1$$

measurement on first qubit \Rightarrow probability

| | |
|-----------------------------|-------------------------------------------------------------------------------------------------------------------|
| suppose result 0 | $\frac{ \alpha_{00} ^2 + \alpha_{01} ^2}{ \alpha_{00} ^2 + \alpha_{01} ^2 + \alpha_{10} ^2 + \alpha_{11} ^2}$ |
| result 1 | $\frac{ \alpha_{10} ^2 + \alpha_{11} ^2}{ \alpha_{00} ^2 + \alpha_{01} ^2 + \alpha_{10} ^2 + \alpha_{11} ^2}$ |

state after measurement:

result 0 $\frac{\alpha_{00}|00\rangle + \alpha_{01}|01\rangle}{\sqrt{|\alpha_{00}|^2 + |\alpha_{01}|^2}}$

result 1 $\frac{\alpha_{10}|10\rangle + \alpha_{11}|11\rangle}{\sqrt{|\alpha_{10}|^2 + |\alpha_{11}|^2}}$

- important qubit state:

Bell state $\Rightarrow \frac{|00\rangle + |11\rangle}{\sqrt{2}}$

(EPR pair)

particularity of qubit state:

suppose measurement on 1st qubit is made \Rightarrow result of measurement on 2nd qubit is then determined with 100% certainty

\rightarrow outcomes are correlated!

outcomes are correlated even if particles
are infinitely apart \Rightarrow what Einstein did not
like about quantum
mechanics

$\Downarrow \Downarrow$
particularity of quantum mechanics

- in general n qubits: $|x_1, \dots, x_n\rangle$

state $\sum_{x_1, \dots, x_n} \alpha(x_1, \dots, x_n) |x_1, \dots, x_n\rangle$

One single n -qubit state stores 2^n bits
of information $x_1 = 0, 1, \dots, x_n = 0, 1$

Quantum Computation

quantum computation: uses quantum gates
in quantum circuits to carry out tasks

Single qubit gates

classical computers: wires, logic gates

single bit classical gates are trivial

no NOT gate $0 \rightarrow 1$
 $1 \rightarrow 0$

quantum NOT gate

$\alpha|0\rangle + \beta|1\rangle \rightarrow \alpha|1\rangle + \beta|0\rangle$

linear transformation

can be written as a matrix $X = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$

state of system, vector $\begin{bmatrix} \alpha \\ \beta \end{bmatrix}$

$$X \begin{bmatrix} \alpha \\ \beta \end{bmatrix} = \begin{bmatrix} \beta \\ \alpha \end{bmatrix}$$

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in general quantum gates acting on single qubits
can be represented as 2×2 matrices

- condition on matrices: must preserve norm!

$$\begin{bmatrix} \alpha' \\ \beta' \end{bmatrix} = \tilde{X} \begin{bmatrix} \alpha \\ \beta \end{bmatrix}$$

$$|\alpha'|^2 + |\beta'|^2 = |\alpha|^2 + |\beta|^2$$

→ matrix which preserves norm: unitary

matrix

$$\tilde{X}^\dagger \tilde{X} = I$$

⇔

$$\langle \alpha' \beta' | \begin{bmatrix} \alpha' \\ \beta' \end{bmatrix} = \langle \alpha \beta | \begin{bmatrix} \alpha \\ \beta \end{bmatrix}$$

⇔

$$\langle \alpha \beta | \underbrace{\tilde{X}^\dagger \tilde{X}} \begin{bmatrix} \alpha \\ \beta \end{bmatrix} = \langle \alpha \beta | \begin{bmatrix} \alpha \\ \beta \end{bmatrix}$$

$$\tilde{X}^\dagger \tilde{X} = I$$

NOT gate $\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$

- in classical computing: only NOT gate is
not entirely trivial

- in quantum computing:

only condition is unitarity on 2×2 matrices

$$Z = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix} \quad (Z\text{-gate})$$

$$H = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \quad (\text{Hadamard-gate})$$

action of Hadamard gate:

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examples: $|0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$

$|1\rangle = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$

$H|0\rangle = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \begin{pmatrix} 1 \\ 0 \end{pmatrix} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 \\ 1 \end{pmatrix} \rightarrow \frac{1}{\sqrt{2}} (|0\rangle + |1\rangle)$

$H|1\rangle = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \begin{pmatrix} 0 \\ 1 \end{pmatrix} \Rightarrow \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)$

- general representation of n single qubit quantum

gates:

$$U = e^{i\alpha} \begin{bmatrix} e^{-i\beta/2} & 0 \\ 0 & e^{i\beta/2} \end{bmatrix} \begin{bmatrix} \cos \frac{\theta}{2} & -i \sin \frac{\theta}{2} \\ i \sin \frac{\theta}{2} & \cos \frac{\theta}{2} \end{bmatrix} \times \begin{bmatrix} e^{-i\phi/2} & 0 \\ 0 & e^{i\phi/2} \end{bmatrix}$$

Multiple qubit gates

classical two bit gates: AND, OR, XOR (exclusive or)

NAND, NOR

AND: $00 \rightarrow$

$$\begin{array}{r|l} & 0 & 1 \\ 0 & 0 & 0 \\ 1 & 0 & 1 \end{array}$$

OR:

$$\begin{array}{r|l} & 0 & 1 \\ 0 & 0 & 1 \\ 1 & 1 & 1 \end{array}$$

XOR:

$$\begin{array}{r|l} & 0 & 1 \\ 0 & 0 & 1 \\ 1 & 1 & 0 \end{array}$$

NAND:

$$\begin{array}{r|l} & 0 & 1 \\ 0 & 1 & 1 \\ 1 & 1 & 0 \end{array}$$

NOR

$$\begin{array}{r|l} & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 0 & 0 \end{array}$$

- NAND - universal gate any operation can be produced as a sequence of NANDs (1)

- example of multi-qubit quantum gate:

controlled-NOT \Rightarrow CNOT

- one-qubit = control qubit

- other-qubit = target qubit

$|A\rangle$ —●— $|A\rangle$ — control

$|B\rangle$ —⊕— $|B \oplus A\rangle$ — target

⊕ - addition modulo 2
(NOT)

$|00\rangle \rightarrow |00\rangle$

$|01\rangle \rightarrow |01\rangle$

$|10\rangle \rightarrow |11\rangle$

$|11\rangle \rightarrow |10\rangle$

\Rightarrow matrix representation

$$\begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

- two qubit gates: 4×4 matrices

constraint: conserve norm-unitarity

- not all classical gates can be generalized to quantum gates

XOR and NAND gates are

not invertible or irreversible

if given output of a classical NAND or XOR \Rightarrow not possible to say what the input was

- quantum gates are always reversible: unitary matrices always have inverses and the inverse is also a unitary matrix
- any multiple qubit gate can be produced from CNOT and single qubit gates

- change of basis: $|+\rangle$

quantum computing can be performed in alternative bases

example:

$$\begin{aligned}
 |+\rangle &= \frac{|0\rangle + |1\rangle}{\sqrt{2}} \\
 |-\rangle &= \frac{|0\rangle - |1\rangle}{\sqrt{2}}
 \end{aligned}
 \Rightarrow
 \begin{aligned}
 \langle + | + \rangle &= \langle - | - \rangle = 1 \\
 \langle + | - \rangle &= \langle - | + \rangle = 0
 \end{aligned}$$

in general $|a\rangle$ & $|b\rangle$

$$\begin{aligned}
 \langle a | a \rangle &= \langle b | b \rangle = 1 \\
 \langle a | b \rangle &= \langle b | a \rangle = 0
 \end{aligned}$$

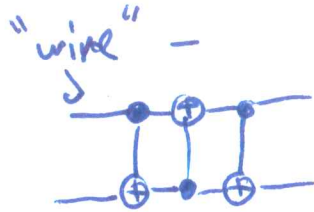
form good bases

rewrite $|0\rangle, |1\rangle$ state in new basis

$$\begin{aligned}
 \alpha |0\rangle + \beta |1\rangle &= \alpha \left(\frac{|+\rangle + |-\rangle}{\sqrt{2}} \right) + \beta \left(\frac{|+\rangle - |-\rangle}{\sqrt{2}} \right) \\
 &= \frac{\alpha + \beta}{\sqrt{2}} |+\rangle + \frac{\alpha - \beta}{\sqrt{2}} |-\rangle
 \end{aligned}$$

Quantum circuits

example:



wire - not necessarily physical wire
can also indicate passage of time (chronology of gates)

what does this circuit do?

$$|a, b\rangle \rightarrow |a, a \oplus b\rangle \rightarrow |a \oplus (a \oplus b), a \oplus b\rangle$$

$$\rightarrow |b, b \oplus (a \oplus b)\rangle$$

$$|b, a\rangle$$

Since $a \oplus (a \oplus b) = b$

proof: \Leftrightarrow
calculate explicitly:

| a | b | $a \oplus b$ | $a \oplus (a \oplus b)$ |
|---|---|--------------|-------------------------|
| 0 | 0 | 0 | 0 |
| 0 | 1 | 1 | 1 |
| 1 | 0 | 1 | 0 |
| 1 | 1 | 0 | 1 |

circuit switches a with b

$$|a, b\rangle \rightarrow |b, a\rangle$$

- in quantum circuits:
- loops not allowed
 - wires can ~~join~~ ^{not join} (FANIN)
 - (OK) \Rightarrow not reversible
 - spreading out of wires (FANOUT)
 - bits copied \Rightarrow not allowed in quantum circuits