

## Chapter 8

# The Prime Number Theorem

The Prime Number Theorem makes predictions about the growth of the prime counting function

$$\pi(x) = \{p \leq x : p \text{ prime}\}.$$

The following diagrams from Zagier's article "The first 50 million primes" (Mathematical Intelligencer (1977), 1-19) show how  $\pi(x)$  behaves for  $x \leq 100$ :

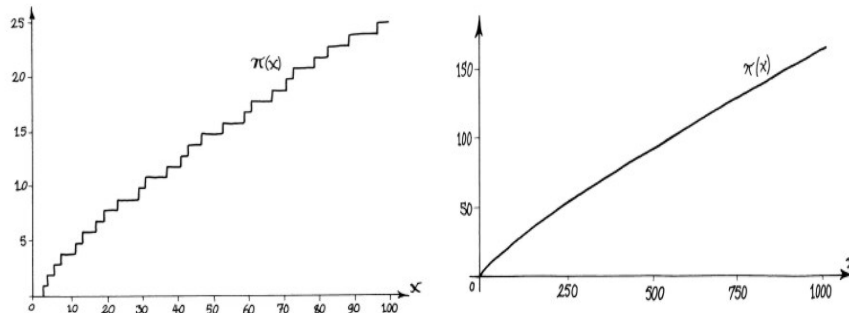


Figure 8.1:  $\pi(x)$  for  $x \leq 100$  and  $x \leq 1000$

It seems clear that  $\pi(x)$  behaves like a smooth function for large values of  $x$ . Is there an analytic function  $f(x)$  such that  $\pi(x) \sim f(x)$ , i.e.,  $\lim \frac{\pi(x)}{f(x)} = 1$ ? Legendre and Gauss conjectures that  $f(x) = \frac{x}{\log x}$  would do it, and this was proved a hundred years later independently by Hadamard and de la Vallée-Poussin using complex analysis (essentially, the truth of the prime number theorem rests upon the fact that the Riemann zeta function defined by  $\zeta(s) = \sum_{n \geq 1} n^{-s}$  in the half plane  $\text{Re } s > 1$  can be extended meromorphically to the whole complex plane and does not have any zeros on the line  $\text{Re } s = 1$ ).

A first step in the direction of a proof of the prime number theorem was done by Chebyshev, who proved with his bare hands that there exist constants  $c_1, c_2 > 0$  with

$$c_1 \frac{x}{\log x} < \pi(x) < c_2 \frac{x}{\log x}.$$

The closer the constants  $c_j$  are to 1, the more technical the proof becomes. Here we will show that  $c_1 = \frac{\log 2}{2}$  and  $c_2 = 6 \log 2$  do work.

Chebyshev's proof uses elementary properties of the binomial coefficients  $\binom{2m}{m} = \frac{(2m)!}{m!m!}$ . Here are some of them:

1.  $\frac{2^{2n}}{2n} \leq \binom{2n}{n} \leq 2^{2n}$ : this comes from  $(1+1)^{2n} = \sum_{m=0}^{2n} \binom{2m}{m}$ .
2.  $\binom{2n}{n}$  is not divisible by any primes  $p > 2n$ .
3.  $\binom{2n}{n}$  is divisible by all primes  $n < p \leq 2n$ .

## 8.1 The Upper Bound

Properties (2) and (1) of the middle binomial coefficient imply that

$$n^{\pi(2n) - \pi(n)} \leq \prod_{n < p \leq 2n} p \leq \binom{2n}{n} \leq 2^{2n}.$$

Taking the log gives  $\pi(2n) - \pi(n) \leq 2 \log 2 \frac{n}{\log n}$ . Using induction we now easily see that

$$\pi(2^k) \leq 3 \cdot \frac{2^k}{k}.$$

In fact, this is checked directly for  $k \leq 5$ ; for  $k > 5$  we find

$$\pi(2^{k+1}) \leq \pi(2^k) + \frac{2^{k+1}}{k} \leq \frac{3 \cdot 2^k}{k} + \frac{2 \cdot 2^k}{k} = \frac{5 \cdot 2^k}{k} \leq \frac{3 \cdot 2^{k+1}}{k+1}.$$

Now we exploit the fact that  $f(x) = \frac{x}{\log x}$  is monotonely increasing for  $x \geq e$ . Thus if  $4 \leq 2^k < x \leq 2^{k+1}$ , then

$$\pi(x) \leq \pi(2^{k+1}) \leq 6 \frac{2^k}{k+1} \leq 6 \log 2 \frac{2^k}{\log 2^k} \leq 6 \log 2 \frac{x}{\log x}.$$

Since  $\pi(x) \leq 6 \log 2 \frac{x}{\log x}$  for  $x \leq 4$ , the proof is now complete.

## 8.2 The Lower Bound

Here we will have to work slightly harder. First we prove

**Lemma 8.1.** *Let  $v_p(n)$  denote the exponent of the maximal power of  $p$  dividing  $n$ . Then*

$$v_p(n!) = \sum_{m \geq 1} \left\lfloor \frac{n}{p^m} \right\rfloor.$$

*Proof.* Among the numbers  $1, 2, \dots, n$ , exactly  $\lfloor \frac{n}{p} \rfloor$  are multiples of  $p$  and thus contribute 1 to the exponent; moreover, exactly  $\lfloor \frac{n}{p^2} \rfloor$  are multiples of  $p^2$  and contribute another 1 to the exponent ...  $\square$

Now put  $N = \binom{2n}{n}$ . By the lemma above we have

$$v_p(N) = \sum_{m \geq 1} \left( \left\lfloor \frac{2n}{p^m} \right\rfloor - 2 \left\lfloor \frac{n}{p^m} \right\rfloor \right).$$

Now we use

**Lemma 8.2.** *For all  $x \in \mathbb{R}$  we have  $\lfloor 2x \rfloor - 2\lfloor x \rfloor \in \{0, 1\}$ .*

*Proof.* Write  $x = \lfloor x \rfloor + \{x\}$ . If the fractional part  $\{x\} < \frac{1}{2}$ , then  $2x = 2\lfloor x \rfloor + \{2x\}$ , hence  $\lfloor 2x \rfloor - 2\lfloor x \rfloor = 0$ . If  $\{x\} \geq \frac{1}{2}$ , then we get  $\lfloor 2x \rfloor - 2\lfloor x \rfloor = 1$ .  $\square$

It is also clear that if  $m > \frac{\log 2n}{\log p}$ , then  $\lfloor \frac{2n}{p^m} \rfloor - 2\lfloor \frac{n}{p^m} \rfloor = 0$ . Thus we find  $v_p(N) \leq \lfloor \frac{\log 2n}{\log p} \rfloor$ , and now

$$\begin{aligned} 2n \log 2 - \log 2n &\leq \log \binom{2n}{n} && \text{because } \frac{2^{2n}}{2n} \leq \binom{2n}{n} \\ &\leq \sum_{p \leq 2n} \left\lfloor \frac{\log 2n}{\log p} \right\rfloor \log p && \text{because } N = \prod p^{v_p(N)} \\ &\leq \sum_{p \leq 2n} \log 2n && \text{because } \lfloor x \rfloor \leq x \\ &= \pi(2n) \log 2n \end{aligned}$$

This yields the lower bound

$$\pi(2n) \geq \log 2 \frac{2n}{\log 2n} - 1.$$

We claim that this implies

$$\pi(x) \geq \frac{\log 2}{2} \frac{x}{\log x}$$

for all  $x \geq 2$ . This inequality can be checked directly for  $x \leq 16$ , hence it is sufficient to prove it for  $x > 16$ . Pick an integer  $n$  with  $16 \leq 2n < x \leq 2n + 2$ . Then

$$\frac{2n}{\log 2n} - \frac{n+1}{\log 2n} = \frac{n-1}{\log 2n} \geq \frac{7}{4 \log 2} \geq \frac{1}{\log 2},$$

hence

$$\pi(x) \geq \pi(2n) \geq \log 2 \frac{2n}{\log 2n} - 1 \geq \frac{(n+1) \log 2}{\log(2n+2)} \geq \frac{\log 2}{2} \frac{x}{\log x}.$$

### 8.3 Bertrand's Postulate

By improving the constants  $c_1$  and  $c_2$  above it is clear that Bertrand's conjecture that  $\pi(2n) - \pi(n) > 0$  for all integers  $n$  must follow: there is at least one prime between  $n$  and  $2n$  for all  $n \in \mathbb{N}$ . In fact a simple calculation shows that Bertrand's conjecture follows as soon as we can find bounds that satisfy  $2c_1 > c_2$ .

Here is a direct proof. Put  $\Theta(x) = \prod_{p \leq x} p$ .

**Proposition 8.3.** *We have  $\Theta(x) \leq 4^x$ .*

*Proof.* It is clearly sufficient to prove this for integers  $x \geq 4$ . Observe that  $\binom{2m+1}{m} = \binom{2m+1}{m+1} < 2^{2m} = 4^m$ . This gives

$$\prod_{m+2 \leq p \leq 2m+1} p \mid \binom{2m+1}{m} < 4^m.$$

Now we prove the claim by induction; assume it is true for all  $n < k$ . If  $k$  is even, then  $\Theta(k) = \Theta(k-1) < 4^{k-1} < 4^k$  by induction assumption and the fact that  $k$  is not prime. Assume therefore that  $k = 2m+1$ . Then

$$\Theta(k) = \Theta(m+1) \prod_{m+2 \leq p \leq 2m+1} p < 4^{m+1} \cdot 4^m = 4^k.$$

□

We have already seen that  $N = \binom{2n}{n}$  is divisible by all primes  $p$  with  $n < p \leq 2n$  (if there are any). Now we claim that the primes  $p$  with  $\frac{2}{3}n < p \leq n$  do not divide  $N$ . In fact we have  $2n < 3p \leq p^2$ , hence  $2 \leq \frac{2n}{p} < 3$ , hence  $\lceil \frac{2n}{p} \rceil = 2$  and  $\lfloor \frac{n}{p} \rfloor = 1$ , and this implies that  $v_p(N) = 2 - 2 = 0$ .

Now we prove Bertrand's postulate by contradiction. Assume there is an integer  $n$  such that the interval  $(n, 2n]$  does not contain any prime. By the discussion above this implies that  $N = \binom{2n}{n}$  is not divisible by any prime  $p > \frac{2}{3}n$ . Now consider primes  $p \mid N$  with  $v_p(N) > 1$ . They satisfy  $p^2 \leq p^{v_p} \leq 2n$ , hence we must have  $p \leq \sqrt{2n}$  for such primes. The number of such primes is clearly bounded by  $\sqrt{2n}$ . Now we find

$$\frac{2^{2n}}{2^n} \leq \binom{2n}{n} \leq \prod_{v_p > 1} p^{v_p} \cdot \prod_{v_p = 1} p \leq (2n)^{\sqrt{2n}} \cdot \Theta\left(\frac{2n}{3}\right) \leq (2n)^{\sqrt{2n}} 2^{4n/3}.$$

Taking the log we get

$$2n \log 2 \leq 3(1 + \sqrt{2n}) \log 2n.$$

Since  $\frac{x}{\log x}$  is monotonically increasing for  $x > 3$ , this inequality must be false for all sufficiently large values of  $n$ . In fact, it is false for  $n \geq 512$ . For  $n < 512$ , Bertrand's postulate is proved by looking at the sequence of primes 7, 13, 23, 43, 83, 163, 317, 631.